# FLINT 1.0: Fast Library for Number Theory

William B. Hart and David Harvey

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# 1 Introduction

FLINT is a C library of functions for doing number theory. It is highly optimised and can be compiled on numerous platforms. FLINT also has the aim of providing support for multicore and multiprocessor computer architectures, though we do not yet provide this facility.

FLINT is currently maintained by William Hart of Warwick University in the UK and David Harvey of Harvard University in the US.

As of version 1.0, FLINT compiles on and supports 32 and 64 bit x86 processors, the G5 and Alpha processors, though in theory it compiles on any machine with gcc version 3.4 or later and with GMP version 4.2.1 or later.

FLINT is supplied as a set of modules, fmpz, fmpz\_poly, etc., each of which can be linked to a C program making use of their functionality.

All of the functions in FLINT have a corresponding test function provided in an appropriately named test file, e.g. all the functions in the file fmpz\_poly.c have test functions in the file fmpz\_poly-test.c.

# 2 Building and using FLINT

The easiest way to use FLINT is to build a shared library. Simply download the FLINT tarball and untar it on your system.

Next, set the environment variables  $FLINT\_GMP\_LIB\_DIR$  and  $FLINT\_GMP\_INC\_DIR$  to point to your GMP library and include directories respectively.

Next type:

source flint\_env

in the main directory of the FLINT directory tree.

Finally type:

make library

Move the library file libflint.so, libflint.dll or libflint.dylib (depending on your platform) into your library path and move all the .h files in the main directory of FLINT into your include path.

Now to use FLINT, simply include the appropriate header files for the FLINT modules you wish to use in your C program. Then compile your program, linking against the FLINT library and GMP with the options -lflint -lgmp.

# 3 Test code

Each module of FLINT has an extensive associated test module. We strongly recommend running the test programs before relying on results from FLINT on your system.

To make the test programs, simply type:

```
make test
```

in the main FLINT directory.

The following is a list of the test programs which should be run:

mpn\_extras-test

fmpz\_poly-test

fmpz-test

ZmodF-test

ZmodF\_poly-test

mpz\_poly-test

ZmodF\_mul-test

long\_extras-test

# 4 Reporting bugs

The maintainers wish to be made aware of any and all bugs. Please send an email with your bug report to hart\_wb@yahoo.com.

If possible please include details of your system, version of gcc, version of GMP and precise details of how to replicate the bug.

Note that FLINT needs to be linked against version 4.2.1 or later of GMP and must be compiled with gcc version 3.4 or later.

# 5 Example programs

FLINT comes with a number of example programs to demonstrate current and future FLINT features. To make the example programs, type:

make examples

The current example programs are:

 $delta_qexp$  Compute the first n terms of the delta function, e.g.  $delta_qexp$  1000000 will compute the first one million terms of the q-expansion of delta.

BPTJCubes Implements the algorithm of Beck, Pine, Tarrant and Jensen for finding solutions to the equation  $x^3 + y^3 + z^3 = k$ .

bernoulli Compute bernoulli numbers modulo a large number of primes.

expmod Computes a very large modular exponentiation.

# 6 FLINT macros

In the file flint.h are various useful macros.

The macro constant FLINT\_BITS is set at compile time to be the number of bits per limb on the machine. FLINT requires it to be either 32 or 64 bits. Other architectures are not currently supported.

```
FLINT_MBS(x) returns the absolute value of a long x. FLINT_MIN(x, y) returns the minimum of two long or two unsigned long values x and y. FLINT_MAX(x, y) returns the maximum of two long or two unsigned long values x and y. FLINT_BIT_COUNT(x) returns the number of binary bits required to represent an unsigned long x.
```

# 7 The fmpz\_poly module

The fmpz\_poly\_t data type represents elements of  $\mathbb{Z}[x]$ . The fmpz\_poly module provides routines for memory management, basic arithmetic, and conversions to/from other types.

Each coefficient of an fmpz\_poly\_t is an integer of the FLINT fmpz\_t type. Each coefficient of an fmpz\_poly\_t has the same number of limbs allocated for it, thus fmpz\_poly\_t polynomials are useful for dense polynomial arithmetic where the coefficients are not wildly different sizes.

Unless otherwise specified, all functions in this section permit aliasing between their input and output arguments.

# 7.1 Simple example

The following example computes the square of the polynomial  $5x^3 - 1$ .

```
#include "fmpz_poly.h"
....
fmpz_poly_t x, y;
fmpz_poly_init(x);
fmpz_poly_init(y);
fmpz_poly_set_coeff_ui(x, 3, 5);
fmpz_poly_set_coeff_si(x, 0, -1);
fmpz_poly_mul(y, x, x);
fmpz_poly_print(x); printf("\n");
fmpz_poly_print(y); printf("\n");
fmpz_poly_clear(x);
fmpz_poly_clear(y);

The output is:
4   -1 0 0 5
7 1 0 0 -10 0 0 25
```

## 7.2 Definition of the fmpz polynomial type

The fmpz\_poly\_t type is a typedef for an array of length 1 of fmpz\_poly\_struct's. This permits passing parameters of type fmpz\_poly\_t 'by reference' in a manner similar to the way GMP integers of type mpz\_t can be passed by reference.

In reality one never deals directly with the struct and simply deals with objects of type fmpz\_poly\_t. For simplicity we will think of an fmpz\_poly\_t as a struct, though in practice to access fields of this struct, one needs to dereference first, e.g. to access the limbs field of an fmpz\_poly\_t called poly1 one writes poly1->limbs.

With this way of thinking, fmpz\_poly\_t then has four fields:

- mp\_limb\_t\* coeffs. This array contains all the fmpz\_t's representing the coefficients of the polynomial, consecutively. The first coefficient represents the constant coefficient of the polynomial.
- unsigned long limbs. The number of limbs allocated for the absolute value of each coefficient. An additional limb per coefficient is also allocated to store the sign/size of the coefficient.
- unsigned long alloc. The maximum number of coefficients which can be stored in coeffs. The total amount of space allocated in coeffs is thus alloc\*(limbs+1).
- unsigned long length. The current length of the polynomial, i.e. the number of coefficients which contain actual data. Always length <= alloc. The polynomial is the zero polynomial if and only if length == 0.

An fmpz\_poly\_t is said to be *normalised* if either length == 0, or if the final coefficient is nonzero. All fmpz\_poly functions expect their inputs to be normalised, and unless otherwise specified they produce output that is normalised.

It is permissible to access the coefficients directly by modifying the limbs in coeffs, however if you modify the coefficients in this way, you must ensure that the polynomial is subsequently normalised by calling fmpz\_poly\_normalise().

## 7.3 Managed versus unmanaged layer

The module fmpz\_poly has two layers, a 'managed' and an 'unmanaged' layer. Functions in the unmanaged layer are differentiated by having a leading underscore, e.g. \_fmpz\_poly\_add.

Functions in the managed layer do all the memory management for the user. One does not need to specify the maximum length or number of limbs per coefficient in advance before using a polynomial object. FLINT reallocates space automatically as the computation proceeds, if more space is required.

As a result of the possible need to reallocate, polynomials modified by functions in the managed layer must have been allocated using the FLINT heap memory manager, i.e. only functions such as fmpz\_poly\_init, without the leading underscore, can be used to allocate polynomials for use as *outputs* of managed functions.

On the other hand, the unmanaged layer does no memory management for the user. Each polynomial must have its coefficient limb size and maximum length set in advance. Both the memory management functions in the unmanaged and the managed layer can be used to this end. In particular the unmanaged layer offers stack based memory management options, though note that no reallocation can occur if this option is used.

A final benefit of the unmanaged layer is that one can operate on a range of coefficients of a polynomial. Functions are provided for attaching an fmpz\_poly\_t object to part of an existing polynomial and acting

on that part, as though it were a separate polynomial. This can avoid making unnecessary copies of data and increase the performance of code.

Some functions are available in either the managed or unmanaged layer but not in both.

We now describe the functions available in fmpz\_poly.

## 7.4 Initialisation and memory management

```
void fmpz_poly_init(fmpz_poly_t poly)
```

Initialise an fmpz\_poly\_t for use. All the fields alloc, length and limbs of poly are set to zero. A corresponding call to fmpz\_poly\_clear must be made after finishing with the fmpz\_poly\_t to free the memory used by the polynomial.

For efficiency reasons, a call to fmpz\_poly\_init does not actually allocate any memory for coefficients. Each of the managed functions will automatically allocate any space needed for coefficients and in fact the easiest way to use the managed layer is to let FLINT do all the allocation automatically.

To this end, a user need only ever make calls to the fmpz\_poly\_init and fmpz\_poly\_clear memory management functions if they so wish. Naturally, more efficient code may result if the other memory management functions are also used.

Initialise an fmpz\_poly\_t for use, allocating space for alloc coefficients each with the given number of limbs of space (plus an additional limb for the sign/size limb for each coefficient). The length field is set to zero.

This function should be used when the maximum length of the polynomial and the size of the coefficients is roughly known in advance. It may be faster than having FLINT automatically increase the size of the polynomial as the computation proceeds.

```
void fmpz_poly_realloc(fmpz_poly_t poly, unsigned long alloc)
```

Shrink or expand the polynomial so that it has space for precisely alloc coefficients. If alloc is less than the current length, the polynomial is truncated (and then normalised), otherwise the coefficients and current length remain unaffected.

If the parameter alloc is zero, any space currently allocated for coefficients in poly is freed. A subsequent call to fmpz\_poly\_clear is still permitted and does nothing.

For performance reasons, if poly->limbs is currently zero, this function does not do any allocation. A subsequent call to fmpz\_poly\_fit\_limbs will do the actual allocation.

```
void fmpz_poly_fit_length(fmpz_poly_t poly, unsigned long alloc)
```

Expand the polynomial (if necessary) so that it has space for at least alloc coefficients. This function will never shrink the memory allocated for coefficients and the contents of the existing coefficients and the current length remain unaffected.

If the limbs field of poly is currently zero, then for performance reasons this function does not actually allocate any space. A subsequent call to fmpz\_poly\_fit\_limbs will do any actual allocation.

```
void fmpz_poly_resize_limbs(fmpz_poly_t poly, unsigned long limbs)
```

Shrink or expand the coefficients so that each of them has space for the given number of limbs. It is required that either the existing coefficients still fit into the new limb size or the parameter limbs is set to zero. Given the former, the contents of the existing coefficients and the current length will remain unaffected.

If the parameter limbs is zero then any space currently allocated for coefficients in poly is freed. A subsequent call to fmpz\_poly\_clear is still permitted and does nothing.

If poly->alloc is currently zero, this function does no allocation.

```
void fmpz_poly_fit_limbs(fmpz_poly_t poly, unsigned long limbs)
```

Expand (if necessary) the coefficients so that each of them has space for the given number of limbs. This function will never shrink coefficients, thus existing coefficients and the current length are always preserved.

For performance reasons, no space is allocated if poly->alloc is currently zero. A subsequent call to fmpz\_poly\_fit\_length will do the actual allocation.

```
void fmpz_poly_clear(fmpz_poly_t poly)
```

Free all memory used by the coefficients of poly. The polynomial object poly cannot be used again until a subsequent call to an initialisation function is made.

Initialise a polynomial, allocating space for alloc coefficients each taking no more than the given number of limbs of space (plus an additional limb for the sign/size).

Space is allocated on the FLINT stack, and can only be released by a corresponding call to \_fmpz\_poly\_stack\_clear.

Polynomials initialised and cleared in this way can only be used by unmanaged functions (with a leading underscore) and as *inputs* to managed functions. Reallocation or changing the number of limbs per coefficient is not permitted.

The alloc and limbs parameters of this function may be zero, in which case no memory is actually allocated. A subsequent call to \_fmpz\_poly\_stack\_clear is still permitted, but does nothing.

```
void _fmpz_poly_stack_clear(fmpz_poly_t poly)
```

Release any space allocated for poly back to the stack. The stack based memory manager requires that polynomials be released in the opposite order to that in which they were initialised with \_fmpz\_poly\_stack\_init.

## 7.5 Setting/retrieving coefficients

Retrieve coefficient n as an mpz\_t.

Coefficients are numbered from zero, starting with the constant coefficient.

The managed version returns zero when  $n \ge poly-\$ length.

Set coefficient n to the value of the given mpz\_t.

Coefficients are numbered from zero, starting with the constant coefficient. If n represents a coefficient beyond the current length of poly, zero coefficients are added in between the existing coefficients and the new coefficient, if required.

Retrieve coefficient n as an fmpz\_t.

Coefficients are numbered from zero, starting with the constant coefficient.

It is assumed that the fmpz\_t supplied has already been allocated with sufficient space for the coefficient being retrieved.

The managed version returns zero when n >= poly->length.

Set coefficient n to the value of the given fmpz\_t.

Coefficients are numbered from zero, starting with the constant coefficient. If n represents a coefficient beyond the current length of poly, zero coefficients are added in between the existing coefficients and the new coefficient, if required.

Return the absolute value of coefficient n as an unsigned long.

Coefficients are numbered from zero, starting with the constant coefficient. If the coefficient is longer than a single limb, the first limb is returned.

The managed version returns zero when  $n \ge poly->length$ .

Set coefficient n to the value of the given unsigned long.

Coefficients are numbered from zero, starting with the constant coefficient. If n represents a coefficient beyond the current length of poly, zero coefficients are added in between the existing coefficients and the new coefficient, if required.

```
long fmpz_poly_get_coeff_si(const fmpz_poly_t poly, unsigned long n)
long _fmpz_poly_get_coeff_si(const fmpz_poly_t poly, unsigned long n)
```

Return the value of coefficient n as a long.

Coefficients are numbered from zero, starting with the constant coefficient. If the coefficient will not fit into a long, i.e. if its absolute value takes up more than FLINT\_BITS - 1 bits then the result is undefined.

The managed version returns zero when  $n \ge poly->length$ .

```
void fmpz_poly_set_coeff_si(fmpz_poly_t poly, unsigned long n, long x)
void _fmpz_poly_set_coeff_si(fmpz_poly_t poly, unsigned long n, long x)
```

Set coefficient n to the value of the given long.

Coefficients are numbered from zero, starting with the constant coefficient. If n represents a coefficient beyond the current length of poly, zero coefficients are added in between the existing coefficients and the new coefficient, if required.

```
fmpz_t fmpz_poly_get_coeff_ptr(fmpz_poly_t poly, unsigned long n)
fmpz_t _fmpz_poly_get_coeff_ptr(fmpz_poly_t poly, unsigned long n)
```

Return a pointer to coefficient n of poly, cast as an fmpz\_t. This function is provided so that individual coefficients can be operated on by functions in the fmpz module.

Coefficients are numbered from zero, starting with the constant coefficient.

The managed version returns NULL when  $n \ge poly->length$ .

# 7.6 String conversions and I/O

The functions in this section are not intended to be particularly fast. They are intended mainly as a debugging aid.

All of the functions use the same string representation of polynomials. It is given by a sequence of integers, in decimal notation, separated by whitespace. The first integer gives the length of the polynomial; the remaining length integers are the coefficients. For example  $5x^3 - x + 1$  is represented by the string "4 1 -1 0 5", and the zero polynomial is represented by "0". The coefficients may be signed and arbitrary precision (provided they fit in the given polynomial).

```
int fmpz_poly_from_string(fmpz_poly_t poly, const char* s)
```

Import a polynomial from a string. If the string represents a valid polynomial the function returns 1, otherwise it returns 0.

```
char* fmpz_poly_to_string(const fmpz_poly_t poly)
```

Convert a polynomial to a string and return a pointer to the string. Space is allocated for the string by this function and must be freed when it is no longer used, by a call to free.

```
void fmpz_poly_fprint(const fmpz_poly_t poly, FILE* f)
```

Convert a polynomial to a string and write it to the given stream.

```
void fmpz_poly_print(const fmpz_poly_t poly)
```

Convert a polynomial to a string and write it to stdout.

```
void fmpz_poly_fread(fmpz_poly_t poly, FILE* f)
```

Read a polynomial from the given stream. Return 1 if the data from the stream represented a valid polynomial, otherwise return 0.

```
void fmpz_poly_read(fmpz_poly_t poly)
```

Read a polynomial from stdin. Return 1 if the data read from stdin represented a valid polynomial, otherwise return 0.

## 7.7 Polynomial parameters (length, degree, limbs, etc.)

```
long fmpz_poly_degree(const fmpz_poly_t poly)
long _fmpz_poly_degree(const fmpz_poly_t poly)
```

Return poly->length - 1. The zero polynomial is defined to have degree -1.

```
unsigned long fmpz_poly_length(const fmpz_poly_t poly) unsigned long _fmpz_poly_length(const fmpz_poly_t poly)
```

Return poly->length. The zero polynomial is defined to have length 0.

```
unsigned long fmpz_poly_limbs(const fmpz_poly_t poly)
unsigned long _fmpz_poly_limbs(const fmpz_poly_t poly)
```

Return poly->limbs.

Each coefficient of poly is allowed up to this many limbs to store its absolute value, plus an additional limb to store its sign/size. Thus the total memory currently allocated for the storage of coefficients is poly->alloc\*(poly->limbs+1).

```
unsigned long fmpz_poly_max_limbs(const fmpz_poly_t poly) unsigned long _fmpz_poly_max_limbs(const fmpz_poly_t poly)
```

Returns the maximum number of limbs required to store the absolute value of coefficients of poly. This may be less than poly->limbs.

```
long fmpz_poly_max_bits(const fmpz_poly_t poly)
long _fmpz_poly_max_bits(const fmpz_poly_t poly)
```

Computes the maximum number of bits b required to store the absolute value of coefficients of poly. If all the coefficients of poly are non-negative, b is returned, otherwise -b is returned.

```
long fmpz_poly_max_bits1(const fmpz_poly_t poly)
long _fmpz_poly_max_bits1(const fmpz_poly_t poly)
```

Computes the maximum number of bits b required to store the absolute value of coefficients of poly. If all the coefficients of poly are non-negative, b is returned, otherwise -b is returned.

The assumption is made that the absolute value of each coefficient fits into an unsigned long. This function will be more efficient than the more general fmpz\_poly\_max\_bits in this situation.

## 7.8 Assignment and basic manipulation

```
void fmpz_poly_set(fmpz_poly_t poly1, const fmpz_poly_t poly2)
void _fmpz_poly_set(fmpz_poly_t poly1, const fmpz_poly_t poly2)
Set polynomial x equal to the polynomial y.

void fmpz_poly_swap(fmpz_poly_t poly1, fmpz_poly_t poly2)
void _fmpz_poly_swap(fmpz_poly_t poly1, fmpz_poly_t poly2)
```

Efficiently swap two polynomials. The coefficients are not moved in memory, pointers are simply switched. The unmanaged version does not swap the alloc fields of the polynomials.

```
void fmpz_poly_zero(fmpz_poly_t poly)
void _fmpz_poly_zero(fmpz_poly_t poly)
```

Set the polynomial to the zero polynomial.

```
void fmpz_poly_zero_coeffs(fmpz_poly_t poly, unsigned long n)
void _fmpz_poly_zero_coeffs(fmpz_poly_t poly, unsigned long n)
```

Set the first n coefficients of poly to zero.

The unmanaged versiom of this function requires that poly have space allocated for at least n coefficients.

```
void fmpz_poly_neg(fmpz_poly_t poly)
void _fmpz_poly_neg(fmpz_poly_t poly)
```

Negate the polynomial, i.e. set it to -poly.

```
void fmpz_poly_truncate(fmpz_poly_t poly, const unsigned long trunc)
void _fmpz_poly_truncate(fmpz_poly_t poly, const unsigned long trunc)
```

If trunc is less than the current length of the polynomial, truncate the polynomial to that length. Note that as the function normalises its output, the eventual length of the polynomial may be less than trunc.

This function considers the polynomial poly to be of length n, notionally truncating and zero padding if required, and reverses the result. Since this function normalises its result the eventual length of output may be less than length.

The unmanaged version of this function requires that output have space allocated for at least length coefficients and that output->limbs is at least poly->limbs.

# 7.9 Subpolynomials

A number of functions are provided for attaching an fmpz\_poly\_t object to an existing polynomial or to a range of coefficients of an existing polynomial providing an alias for the original polynomial or part thereof.

Each of the functions in this section normalise the aliases.

One must take care when manipulating the alias, since manipulating it may leave the original polynomial unnormalised.

One must also be careful that one does not pass a polynomial and an alias for that polynomial to the same function since that function will have no way to tell it is dealing with aliases of the same polynomial.

```
void _fmpz_poly_attach(fmpz_poly_t output, const fmpz_poly_t poly)
```

Attach the fmpz\_poly\_t object output to the polynomial poly. Any changes made to the length field of output then do not affect poly.

Attach the fmpz\_poly\_t object output to poly but shifted to the left by n coefficients. This is equivalent to notionally shifting the original polynomial right (dividing by  $x^n$ ) then attaching to the result.

Attach the  $fmpz_poly_t$  object output to the first n coefficients of the polynomial poly. This is equivalent to notionally truncating the original polynomial to n coefficients then attaching to the result.

```
void _fmpz_poly_normalise(fmpz_poly_t poly)
```

Normalise the polynomial so that either the polynomial is the zero polynomial or the leading coefficient is not zero.

Since all other functions in fmpz\_poly assume that input and output polynomials are normalised, this function is only used when manipulating the internals of a polynomial directly or when using subpolynomials.

# 7.10 Comparison

```
int fmpz_poly_equal(const fmpz_poly_t poly1, const fmpz_poly_t poly2)
int _fmpz_poly_equal(const fmpz_poly_t poly1, const fmpz_poly_t poly2)
```

Return 1 if the two polynomials are equal, 0 otherwise.

## 7.11 Shifting

Shift poly to the left by n coefficients (multiply by  $x^n$ ) and write the result to output. Zero coefficients are inserted.

The unmanaged version of this function requires that output have space allocated for at least n + poly->length coefficients.

The parameter n must be non-negative, but can be zero.

Shift poly to the right by n coefficients (divide by  $x^n$  and discard the remainder) and write the result to output.

The parameter n must be non-negative, but can be zero. Shifting right by more than the current length of the polynomial results in the zero polynomial.

#### 7.12 Addition/subtraction

Set the output to the sum of the input polynomials.

Note that if poly1 and poly2 have the same length, cancellation may occur (if the leading coefficients have the same absolute values but opposite signs) and so the result may have less coefficients than either of the inputs. However, the unmanaged version of this function requires that the output have space allocated for the number of coefficients of the longest of the input polynomials.

When using the unmanaged version, note that overflow may occur when adding coefficients together and so one additional bit may be required to store the output coefficients than was required in either of the input polynomials. The additional bit is only required in the case that overflow occurs.

Set the output to poly1 - poly2.

Note that if poly1 and poly2 have the same length, cancellation may occur (if the leading coefficients have the same values) and so the result may have less coefficients than either of the inputs. However, the unmanaged version of this function requires that the output have space allocated for the number of coefficients of the longest of the input polynomials.

When using the unmanaged version, note that overflow may occur when subtracting coefficients of opposite signs and so one additional bit may be required to store the output coefficients than was required in either of the input polynomials. The additional bit is only required in the case that overflow occurs.

# 7.13 Scalar multiplication and division

Multiply poly by the unsigned long x and write the result to output.

When using the unmanaged version, the coefficients of output must have space for the largest output coefficient, i.e. the sum of the number of bits of the absolute values of x and the largest coefficient of poly.

Multiply poly by the long x and write the result to output.

When using the unmanaged version, the coefficients of output must have space for the largest output coefficient, i.e. the sum of the number of bits of the absolute values of x and the largest coefficient of poly.

Multiply poly by the fmpz\_t x and write the result to output.

When using the unmanaged version, the coefficients of output must have space for the largest output coefficient, i.e. the sum of the number of bits of the absolute values of x and the largest coefficient of poly.

Multiply poly by the mpz\_t x and write the result to output.

Divide poly by the unsigned long x, round quotients towards minus infinity, discard remainders and write the result to output.

When using the unmanaged version, the coefficients of output must have space for the largest input coefficient.

Divide poly by the long x, round quotients towards minus infinity, discard remainders and write the result to output.

When using the unmanaged version, the coefficients of output must have space for the largest input coefficient.

Divide poly by the unsigned long x, round quotients towards zero, discard remainders and write the result to output.

When using the unmanaged version, the coefficients of output must have space for the largest input coefficient.

Divide poly by the long x, round quotients towards zero, discard remainders and write the result to output.

When using the unmanaged version, the coefficients of **output** must have space for the largest input coefficient.

Divide poly by the unsigned long x. Division is assumed to be exact and the result is undefined otherwise.

When using the unmanaged version, the coefficients of output must have space for the largest input coefficient.

Divide poly by the long x. Division is assumed to be exact and the result is undefined otherwise.

When using the unmanaged version, the coefficients of output must have space for the largest input coefficient.

Divide poly by the fmpz\_t x, round quotients towards minus infinity, discard remainders, and write the result to output.

When using the unmanaged version, the coefficients of the polynomial output must have sufficient space allocated for limbs1 - limbs2 + 1 limbs, where limbs1 is the maximum number of limbs of the coefficients in poly and limbs2 is the number of limbs required to store the absolute value of x.

Divide poly by the  $mpz_t$  x, round quotients towards minus infinity, discard remainders, and write the result to output.

#### 7.14 Polynomial multiplication

Multiply the two given polynomials and return the result in output.

When using the unmanaged version, the coefficients of the output polynomial may be as large as bits1 + bits2 + bits(length2) where bits1 is the number of bits of the absolute value of the largest coefficient of poly1, bits2 is the corresponding thing for poly2, bits(length2) is the number of bits in the binary representation of the length of the shortest polynomial.

The length of the output polynomial will be poly1->length + poly2->length - 1.

Multiply the two given polynomials and truncate the result to n coefficients, storing the result in output. This is sometimes known as a short product.

See \_fmpz\_poly\_mul for a discussion of how big the output coefficients can be.

The length of the output polynomial will be at most the minimum of n and the value poly1->length + poly2->length - 1. It is permissible to set n to any non-negative value, however the function is optimised for n about half of poly1->length + poly2->length.

This function is more efficient than multiplying the two polynomials then truncating. It is the operation used when multiplying power series.

Multiply the two given polynomials storing the result in output. This function guarantees all the coefficients except the first n, which may be arbitrary. This is sometimes known as an opposite short product.

See \_fmpz\_poly\_mul for a discussion of how big the output coefficients can be.

The length of the output polynomial will be poly1->length + poly2->length - 1 unless n is greater than or equal to this value, in which case it will return the zero polynomial. It is permissible to set n to any non-negative value, however the function is optimised for n about half of poly1->length + poly2->length.

For short polynomials, this function is more efficient than computing the full product.

# 7.15 Polynomial division

Performs division with remainder in  $\mathbb{Z}[x]$ . Computes polynomials Q and R in  $\mathbb{Z}[x]$  such that the equation A = B\*Q + R, holds. All but the final B->length - 1 coefficients of R will be positive and less than the absolute value of the lead coefficient of R.

Note that in the special cases where the leading coefficient of B is  $\pm 1$  or A = B\*Q for some polynomial Q, the result of this function is the same as if the computation had been done over  $\mathbb{Q}$ .

Performs division without remainder in  $\mathbb{Z}[x]$ . The computation returns the same result as fmpz\_poly\_divrem, but no remainder is computed. This is in general faster than computing quotient and remainder.

Note that in the special cases where the leading coefficient of B is  $\pm 1$  or A = B\*Q for some polynomial Q, the result of this function is the same as if the computation had been done over  $\mathbb{Q}$ . In particular it can be used efficiently for exact division in  $\mathbb{Z}[x]$ .

Performs power series division in  $\mathbb{Z}[[x]]$ . The function considers the polynomials A and B to be power series of length n starting with the constant terms. The function assumes that B is normalised, i.e. that the constant coefficient is  $\pm 1$ . The result is truncated to length n regardless of the inputs.

#### 7.16 Pseudo division

Performs division with remainder of two polynomials in  $\mathbb{Z}[x]$ , notionally returning the results in  $\mathbb{Q}[x]$  (actually in  $\mathbb{Z}[x]$  with a single common denominator).

Computes polynomials Q and R such that lead(B)^d\*A = B\*Q + R where R has degree less than that of B.

This function may be used to do division of polynomials in  $\mathbb{Q}[x]$  as follows. Suppose polynomials  $\mathbb{C}$  and  $\mathbb{D}$  are given in  $\mathbb{Q}[x]$ .

- 1) Write C = d1\*A and D = d2\*B for some polynomials A and B in  $\mathbb{Z}[x]$  and integers d1 and d2.
- 2) Use pseudo-division to compute Q and R in  $\mathbb{Z}[x]$  so that  $1^d*A = B*Q + R$  where 1 is the leading coefficient of B.
- 3) We can now write  $C = (d1/d2*D*Q + d1*R)/1^d$ .

Performs division without remainder of two polynomials in  $\mathbb{Z}[x]$ , notionally returning the results in  $\mathbb{Q}[x]$  (actually in  $\mathbb{Z}[x]$  with a single common denominator).

Notionally computes polynomials Q and R such that  $lead(B)^d*A = B*Q + R$  where R has degree less than that of B, but returns only Q. This is slightly more efficient than computing the quotient and remainder.

## 7.17 Powering

Raises poly to the power exp and writes the result in output.

Notionally raises poly to the power exp, truncates the result to length n and writes the result in output. This is computed much more efficiently than simply powering the polynomial and truncating.

This function can be used to raise power series to a power in an efficient way.

# 8 The fmpz module

The fmpz module is designed for manipulation of the FLINT flat multiprecision integer format fmpz\_t. An fmpz\_t is not a struct but merely a pointer to an array of limbs laid out in a certain way.

The first limb is a sign/size limb. If it is 0 the integer represented by the fmpz\_t is 0. The absolute value of the sign/size limb is the number of subsequent limbs that the absolute value of the integer being represented, takes up. The absolute value of the integer is then stored as limbs, least significant limb first, in the subsequent limbs after the sign/size limb. If the sign/size limb is positive, a positive integer is intended and if the sign/size limb is negative the negative integer with the stored absolute value is intended.

The fmpz\_t type is not intended as a standalone integer type. It is intended to be used in composite types such as polynomials and matrices which consist of many integer entries. All memory management is then done by the composite type, not by the fmpz module itself. Thus, none of the functions in the fmpz module do automatic memory management. It is up to the user to ensure that output fmpz\_t's have sufficient space allocated for them.

## 8.1 A simple example

We start with a simple example of the use of the fmpz module.

This example sets x to 3 and adds 5 to it.

```
#include "fmpz.h"
....

fmpz_t x = fmpz_init(1); // Allocate 1 limb of space
fmpz_set_ui(x, 3);
fmpz_add_ui_inplace(x, 5);
printf("3_+_5_is_"); fmpz_print(x); printf("\n");
fmpz_clear(x);
```

We now discuss the functions available in the fmpz module.

## 8.2 Memory management

```
fmpz_t fmpz_init(unsigned long limbs)
```

Allocates space for an fmpz\_t with the given number of limbs (plus an additional limb for the sign/size) on the heap and return a pointer to the space.

```
fmpz_t fmpz_realloc(fmpz_t f, unsigned long limbs)
```

Reallocate the space used by the fmpz\_t f so that it has space for the given number of limbs (plus a sign/size limb). The parameter limbs must be non-negative. The existing contents of f are not altered if they still fit in the new size.

```
void fmpz_clear(const fmpz_t f)
```

Free space used by the fmpz\_t f.

```
fmpz_t fmpz_stack_init(unsigned long limbs)
```

Allocates space for an fmpz\_t with the given number of limbs (plus an additional limb for the sign/size) on the stack and return a pointer to the space.

```
void fmpz_stack_clear(const fmpz_t f)
```

Return space used by the fmpz\_t f to the stack.

## 8.3 String operations

```
void fmpz_print(const fmpz_t f)
```

Print the multiprecision integer f.

#### 8.4 fmpz properties

```
unsigned long fmpz_size(const fmpz_t f)
```

Return the number of limbs used to store the absolute value of the multiprecision integer f.

```
unsigned long fmpz_bits(const fmpz_t f)
```

Return the number of bits required to store the absolute value of the multiprecision integer f.

```
long fmpz_sgn(const fmpz_t f)
```

Return the sign/size limb of the multiprecision integer f. The sign of the sign/size limb is the sign of the multiprecision integer. The absolute value of the sign/size limb is the size in limbs of the absolute value of the multiprecision integer f.

# 8.5 Assignment

```
void fmpz_set_ui(fmpz_t res, unsigned long x)
    Set the multiprecision integer res to the unsigned long x.
void fmpz_set_si(fmpz_t res, long x)
    Set the multiprecision integer res to the long x.
void fmpz_set(fmpz_t res, const fmpz_t f)
    Set the multiprecision integer res to equal the multiprecision integer f.
8.6
     Comparison
int fmpz_equal(const fmpz_t f1, const fmpz_t f2)
     Return 1 if f1 is equal to f2, otherwise return 0.
int fmpz_is_one(const fmpz_t f)
     Return 1 if f is one, otherwise return 0.
int fmpz_is_zero(const fmpz_t f)
     Return 1 if f is zero, otherwise return 0.
8.7 Conversion
void mpz_to_fmpz(fmpz_t res, const mpz_t x)
    Convert the mpz_t x to the fmpz_t res.
void fmpz_to_mpz(mpz_t res, const fmpz_t f)
    Convert the fmpz_t f to the mpz_t res.
```

# 8.8 Addition/subtraction

```
void fmpz_add(fmpz_t res, const fmpz_t f1, const fmpz_t f2)
    Set res to the sum of f1 and f2.
void fmpz_add_ui_inplace(fmpz_t res, unsigned long x)
    Set res to the sum of res and the unsigned long x.
void fmpz_add_ui(fmpz_t res, const fmpz_t f, unsigned long x)
    Set res to the sum of f and the unsigned long x.
void fmpz_sub(fmpz_t res, const fmpz_t f1, const fmpz_t f2)
    Set res to f1 minus f2.
void fmpz_sub_ui_inplace(fmpz_t res, unsigned long x)
    Set res to res minus the unsigned long x.
void fmpz_sub_ui(fmpz_t res, const fmpz_t f, unsigned long x)
    Set res to f minus the unsigned long x.
8.9 Multiplication
void fmpz_mul(fmpz_t res, const fmpz_t f1, const fmpz_t f2)
    Set res to f1 times f2.
void fmpz_mul_ui(fmpz_t res, const fmpz_t f1, unsigned long x)
    Set res to f1 times the unsigned long x.
void fmpz_addmul(fmpz_t res, const fmpz_t f1, const fmpz_t f2)
    Set res to res + f1 * f2.
```

#### 8.10 Division

```
void fmpz_tdiv(fmpz_t res, const fmpz_t f1, const fmpz_t f2)
```

Set res to the quotient of f1 by f2. Round the quotient towards zero and discard the remainder.

```
void fmpz_fdiv(fmpz_t res, const fmpz_t f1, const fmpz_t f2)
```

Set res to the quotient of f1 by f2. Round the quotient towards minus infinity and discard the remainder.

```
void fmpz_tdiv_ui(fmpz_t res, const fmpz_t f1, unsigned long x)
```

Set **res** to the quotient of **f1** by the unsigned long **x**. Round the quotient towards zero and discard the remainder.

## 8.11 Powering

```
void fmpz_pow_ui(fmpz_t res, const fmpz_t f, unsigned long exp)
```

Set res to f raised to the power exp. This requires exp to be non-negative.

## 8.12 Number theoretical

```
void fmpz_binomial_next(fmpz_t next, const fmpz_t prev, long n, long k)
```

Assuming prev is set to the binomial coefficient bin(n, k-1) this function returns the binomial coefficient bin(n, k). For efficiency reasons, this function requires that next has space for one more limb than the size of prev.

#### 8.13 Miscellaneous

```
void fmpz_normalise(const fmpz_t f)
```

Normalise the multiprecision integer f.

Since all the functions in fmpz assume that all inputs are normalised and all outputs are normalised, this function is usually used internally by FLINT or can be used when modifying the internals of an fmpz\_t.

# 9 The quadratic sieve

Currently the quadratic sieve is a standalone program which can be built by typing:  $\mathtt{make}\ \mathsf{QS}$ 

in the main FLINT directory.

The program is called mpQs. Upon running it, one enters the number to be factored at the prompt.

The quadratic sieve requires that the number entered not be a prime, not be a perfect power and it must not have very small factors. Trial division and the elliptic curve method should be run before making a call to the quadratic sieve, to remove small factors. The sieve may fail silently if the conditions are not met.

# 10 Large integer multiplication

In the module mpn\_extras and mpz\_extras are functions F\_mpn\_mul and F\_mpz\_mul respectively which are drop in replacements for GMP's mpn\_mul and mpz\_mul respectively.

These replacement functions are substantially faster than GMP 4.2.1 when multiplying integers which are thousands of limbs in size. For smaller multiplications these functions call their respective GMP counterparts.